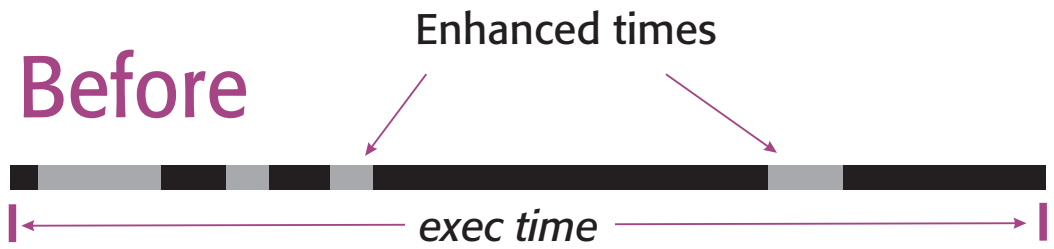
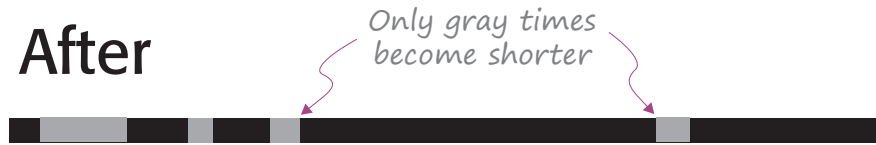


# Enhancing A Part

Parts of a program are enhanced, affected times marked gray.



How much improvement in the execution time should be expected?



Conventionally, put bigger number in numerator to get non-fractional (>1) valued speedup ratio.

$$Speedup_{overall} = \frac{(exec\ time)_{old}}{(exec\ time)_{new}} = ?$$

# Amdahl's Insight

Before

Collect to-be-enhanced times to visualize



Let  $(task\ time)_{old}$  = total times affected by enhancement.

$(task\ time)_{old}$

Only a fraction of the original time will enjoy a full enhanced speedup!

$$Frac_{enh} = \frac{(task\ time)_{old}}{(exec\ time)_{old}}$$

$$Speedup_{enh} = \frac{(task\ time)_{old}}{(task\ time)_{enh}}$$

After

Only grayed parts shrink

Overall speedup will be limited by that fraction of time enhanced,  $Frac_{enh}$

Clearly from the after diagram ...

same unaffected time

$$(exec\ time)_{new} = (exec\ time)_{old} - (task\ time)_{old} + (task\ time)_{enh}$$

Fairly obvious relation leads to a speedup formula after a division and a couple of strategic substitutions (next).

# Amdahl's Law A Formula

$$(exec\ time)_{new} = (exec\ time)_{old} - (task\ time)_{old} + (task\ time)_{enh}$$

$\div (exec\ time)_{old}$   
To obtain an overall speedup term in the left-hand side  
(in inverted form, flip back at the end)

and note

$$\frac{(task\ time)_{old}}{(exec\ time)_{old}} = Frac_{enh}$$

also observe (from previous slide)

another handy substitution for 3rd right-hand term after division

**Exercise**  
Complete the derivation leading to the formula (next). *Ans. last slide.*

$$\frac{Frac_{enh}}{Speedup_{enh}} = \frac{(task\ time)_{old}}{(exec\ time)_{old}} \times \frac{(task\ time)_{enh}}{(task\ time)_{old}}$$

# Amdahl's Law

For example, can't add 2nd processor and expect two-fold improvement in run time unless new processor is usable 100% of that time (i.e.,  $Frac_{enh} = 1$ , which leads to overall speedup of 2).

## Quiz

Evaluate the limit as  $Speedup_{enh} \rightarrow \infty$ . Interpret results.

The limit clearly reveals that the fraction time not enhanced places an upper limit on overall speedup.

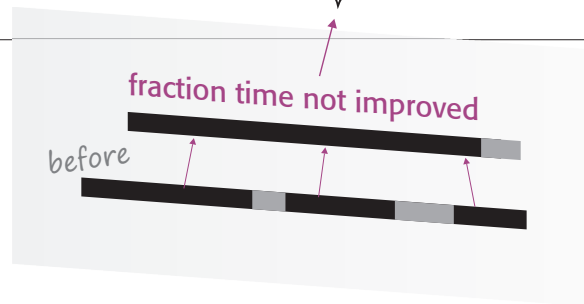


## Exercise

Plot the speedup, overall vs. enhanced (e.g., 2 to 200 times) for different values of  $Frac_{enh}$  (e.g., 0.9, 0.8, ..., 0.4). Comment on the shapes of the curves.

$$\frac{(exec\ time)_{old}}{(exec\ time)_{new}}$$

$$Speedup_{overall} = \frac{1}{(1 - Frac_{enh}) + \frac{Frac_{enh}}{Speedup_{enh}}}$$



# Example 1 (CA:AQA)

**Solution Hint.** Figure out formula variables from problem wording, sometimes you need a small calculation.

Suppose that we are considering an **enhancement that runs 10 times faster than the original machine but is only usable 40% of the time.** What is the overall speedup gained by incorporating the enhancement?

# Example 2 (CA:AQA)

**Suppose a cache is 10 times faster than main memory, and suppose that the cache can be used 90% of the time. How much speedup do we gain by using the cache?**

# Example 3

*Insight builder alert..  
inspect the formula  
carefully first*

**Suppose a program runs in 100 seconds on a computer, with multiply operations responsible for 80 seconds of this time. How much would multiplication need to be improved to make program run five times faster?**

# References & Answers

**Exercise**  
 (Graduate) Check the derivation of Amdahl's law in the original 1967 paper (!)

- Gene M. Amdahl, *Validity of the single processor approach to achieving large scale computing capabilities*, Proceeding AFIPS '67 (Spring) April 18-20, 1967, Spring Joint Computer Conference, pages 483-485.
- [CA:AQA] Hennessy and Patterson, *Computer Architecture: A Quantitative Approach*, Morgan-Kaufmann (ed. 2).

