

Polynomial Evaluation

⇒ Polynomial degree

⇒ Calculate p at an x -value (a point)

$$p(x) = a_n x^n + \cdots + a_1 x + a_0$$

⇒ Seems to involve

 Computing n terms of form $a_i c^i$

 Exponentiation of some constant

⇒ How well can we exponentiate?

Definition naturally suggests multiplication as a basic operation.

Polynomial Evaluation Brute Force Approach

$$2x^4 - x^3 + 3x^2 + x - 5$$

Multiplications, $M(n) = ?$

Quiz
How many multiplications
are needed for the example?

 How many per term? $x^i, a_i \times x^i$

 How many terms? Note decreasing exponent

Quiz
Is this approach favorable for
a **multipoint** (c_1, c_2, \dots, c_m)
evaluation scenario?

 Resulting efficiency

Polynomial Evaluation Representation Change

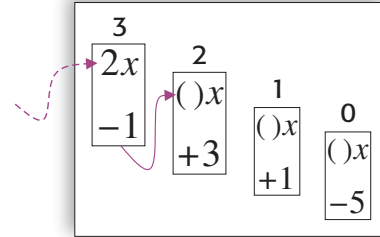
⇨ **Horner's rule**

Exercise
 Transform $p(x)$ to the alternate algebraic form.

$$p(x) = 2x^4 - x^3 + 3x^2 + x - 5$$

Each column (nested factor) may correspond to an iteration in a *for*-loop.

$$= x(x(x(2x - 1) + 3) + 1) - 5$$



$$x^n = xx^{n-1}$$

Once $n-1$ power of x is obtained, do we really need to recompute it to get the next power?

⇨ **Observation**

⇨ **$M(n) = ?$ By inspection? In general (guess)**

Helps to view an initial inner factor associated with coefficient a_n .

⇨ **Pen-paper example ($x=3$)** Focus on developing the procedure rather than getting a final result.

$$x(x(x(2x - 1) \dots$$

Polynomial Evaluation

Horner's Rule

Quiz

What's the efficiency if addition was chosen as basic operation?

Algorithm *Horner*

Input $P[0..n]$ coefficients $a_0 \cdots a_n$ of polynomial p , point x

Output Polynomial value $p(x)$

- 1: $p \leftarrow P[n]$
 - 2: **for** $i \leftarrow n - 1$ **downto** 0 **do**
 - 3: $p \leftarrow x \times p + P[i]$
 - 4: **return** p
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⇒ **Efficiency**

⇒ **Applications**

Representation Change Binary Exponentiation

⇒ Successive squaring

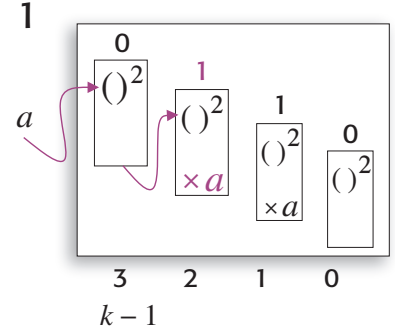
Exercise

Use the diagram to generate a^{22} (binary 10110) where k is #steps. Compare to the diagram depicting polynomial evaluation via Horner's rule.

⇒ **Key idea**

 Successive squaring

 Simple examples



The calculation sequence suggested by Horner's rule + an older idea of exponentiation via successive squaring lead to algorithms that utilize a change of representation of the exponent.

⇒ **Pen-paper procedure**

⇒ **Algorithms (next)**

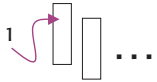
Binary Exponentiation An Algorithm



Exercise

Compare to a decrease-by-constant factor based on the formula $(a^{n/2})^2$. **Hint:** write the recurrence.

Essentially, iterate on an exponent's logarithm (rather than the exponent itself).



Exercise

Modify the pseudocode to initialize p with 1. Will performance change?

Algorithm *LeftRightBinaryExponentiation*

Input Number a

Input Binary representation $b_l \cdots b_1 b_0$ of integer exponent $n > 0$

Output a^n $l \leftarrow k$

- 1: $p \leftarrow a$
- 2: **for** $i \leftarrow l - 1$ **downto** 0 **do**
- 3: $p \leftarrow p \times p$
- 4: **if** $b_i = 1$ **then**
- 5: $p \leftarrow p \times a$
- 6: **return** p

Efficiency?

Representation Change Conclusions



Exercise

Write a recurrence for a divide-conquer exponentiation. Is it a good idea?

Hint: use *WolframAlpha* to check solution.

⇒ **Exponentiation strategies**

⇒ **Polynomial evaluation**



A representation change proves to be a better strategy than trying to improve exponentiation performance.



Via exponentiation



Using Horner Rule

Exercise



Lookup the FFT computation (problem, algorithms, compare efficiencies).

⇒ **Multipoint scenario**

Connections Styles of Iteration



Exercise

Write a summation or a recurrence for definition-based bottom-up and top-down decrease-by-1 algorithms to calculate a^n .

⇒ **Bottom-up: iterative typically**
Build from basic case & work way up

⇒ **Top-down: recursive typically**
Work way down to basic case

⇒ **Example: exponentiation**
Binary vs. decrease-by-const-factor