

Convex Hull Review

⇒ Extreme point

Quiz

Is the set of points $P_1 \dots P_n$ a convex set?

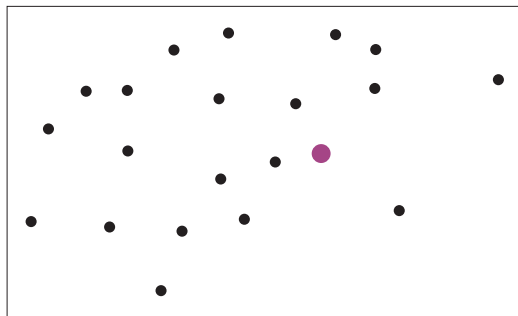
Quiz

Is the highlighted point in the convex hull of the set? Can it be a vertex of an enclosing polygon? Why?

Exercise

Draw the convex hull of the set.

$$P_1 = (x_1, y_1), \dots, P_n = (x_n, y_n)$$



Examine line segments (distinct pairs of points) against each of the remaining $n-2$ points.

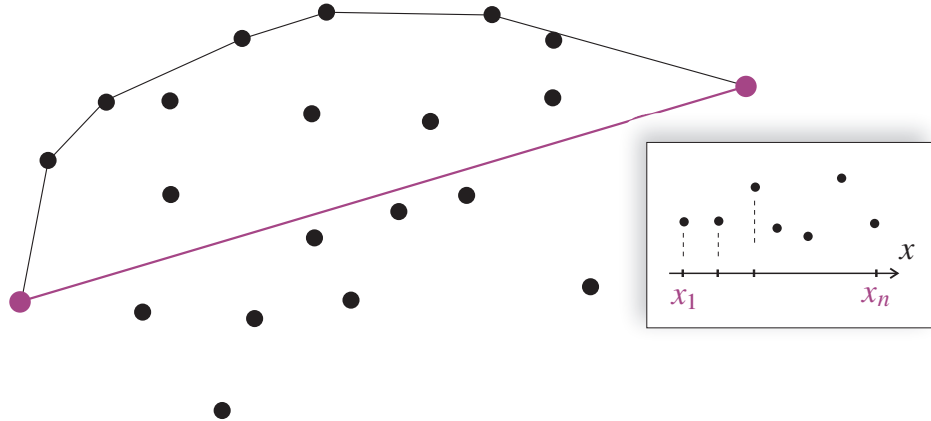
Brute force $O(n^3)$, can we do better?

Divide-Conquer Convex Hull Initializations

⇒ Upper/lower hull

Quiz
How to quickly determine
left-most and right-most
points? What about points
on $\overline{P_1P_n}$?

$$P_1 = (x_1, y_1), \dots, P_n = (x_n, y_n)$$



Divide-Conquer Convex Hull Basic Procedure

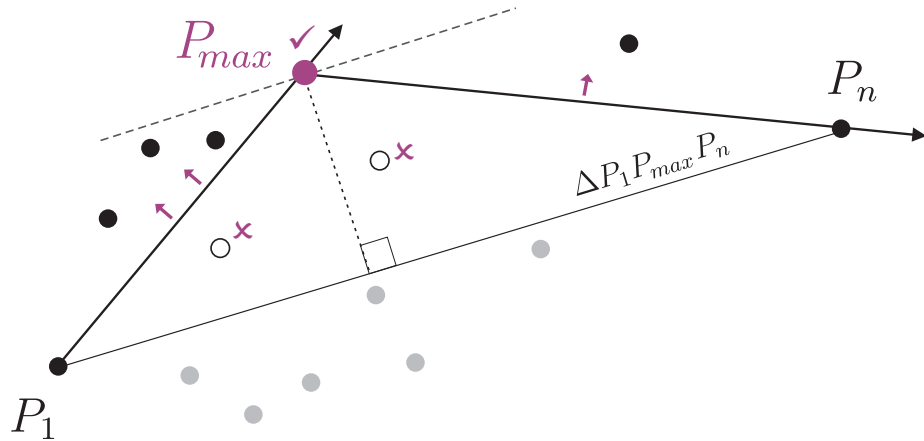


P_{max} divides the upper set into 3 sets.

3 Observations

Quiz

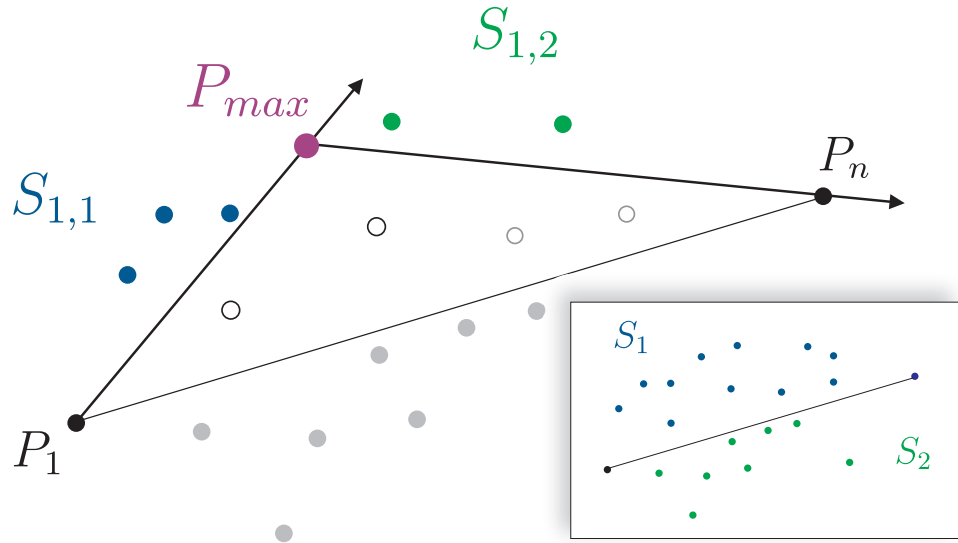
How can the highlighted point be selected? **Hint:** see details slide.



Divide-Conquer Convex Hull A Partition



Each partition identifies a vertex on the upper hull + creates 2 sets to check next.



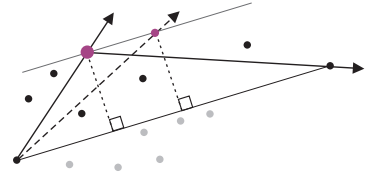
A Quickhull Algorithm Computational Details



Quiz

What if many points qualify to partition set?

⇒ **How to determine P_{max} ?**

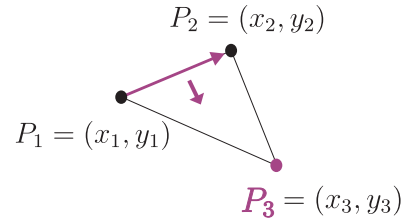


Quiz

What is the sign of the determinant in the figure?

⇒ **Use result for $\Delta P_1 P_2 P_3$**

$$\begin{vmatrix} x_1 & y_1 & 1 \\ x_2 & y_2 & 1 \\ x_3 & y_3 & 1 \end{vmatrix} = x_1 y_2 + x_3 y_1 + x_2 y_3 - x_3 y_2 - x_2 y_1 - x_1 y_3$$

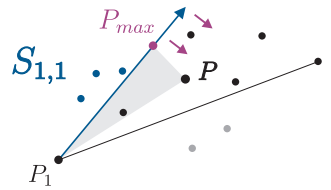


Quiz

Write an expression for $S_{1,2}$.

⇒ **To determine $S_{1,1}$**

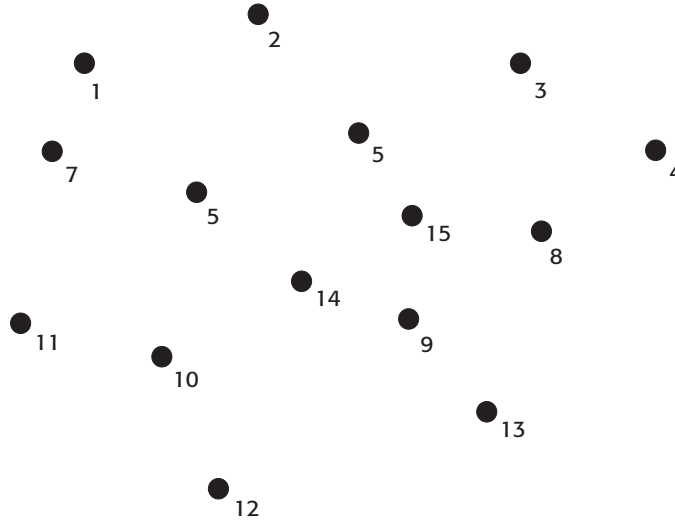
$\forall P = (x, y) \in S_1$, check $\det \Delta \overrightarrow{P_1 P_{max}} P$



A Quickhull Algorithm Exercise



Exercise
Perform a *Quickhull* by hand.
Show for each step: P_{max} ,
excluded points, sets of points to
consider next, and the resulting
convex hull vertices.



Divide-Conquer Closest Pair Basic Procedure

$$S = \{P_1, P_2, P_3, P_4\}$$

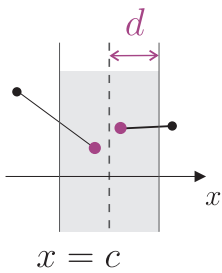
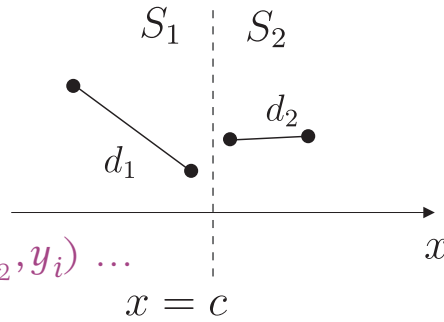
$$(x_1, y_1) \dots$$



Quiz

How to quickly determine a partition point C ? How many points in each subset in general?

Hint: think about algorithm inputs (x_{ik}, y_{ik}) : i original index, k ordered/sorted index).



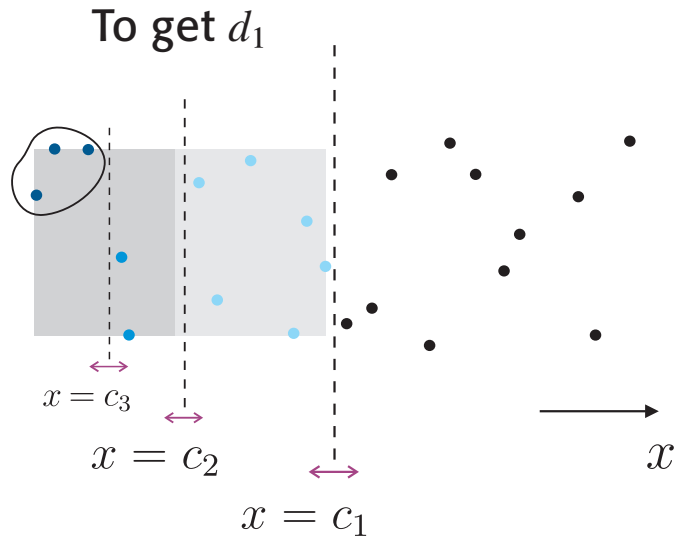
Determine separately, combine

 $d = \min(d_1, d_2)$

 Check pairs around $x = c$

Divide-Conquer Closest Pair Repeat Recursively

Exercise
Design top level pseudocode
without detailing the merge
step.



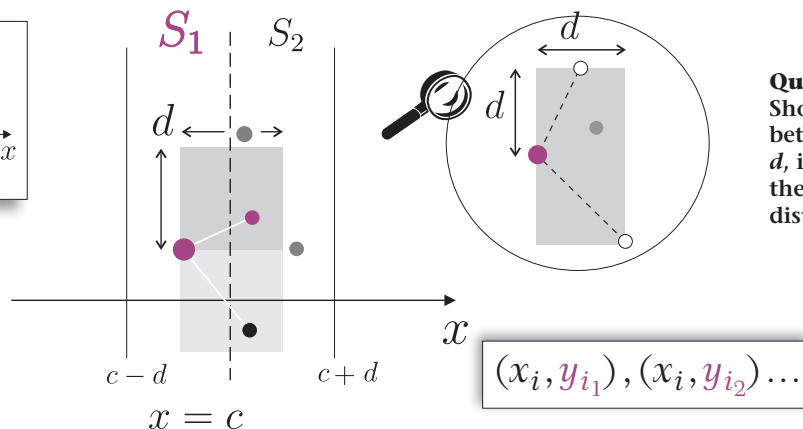
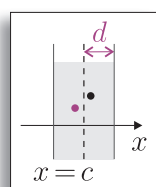
Divide if points > 2 (or 3)
otherwise return a distance.

Keep dividing until set is 3 or less points

Divide-Conquer Closest Pair Merge Procedure

Points outside clearly farther than d but not all inside are closer.

Even inside P_i 's box some points are farther away than d .



Quiz
 Show that the distance between the points is $> d$, if the x-coordinate of the top point is at distance $d/2$.

$$\forall P_i \in S_1(c - d \leq x \leq c), \text{ check } \{P_j \in S_2(c \leq x \leq c + d)\}$$

➡ Major result: no more than six neighbors to consider per point

Performance Results

Quiz
Compare problem size
reduction in each case.



⇒ **Recurrence pattern**

Quiz
Write the recurrence,
determine order of growth
of solution.

⇒ **Divide-conquer closest pair**



⇒ **Quickhull worst and average**

Exercise
Lookup best and average
case efficiencies of the
brute force convex hull of
Chapter 3.

⇒ **Compare to brute force**